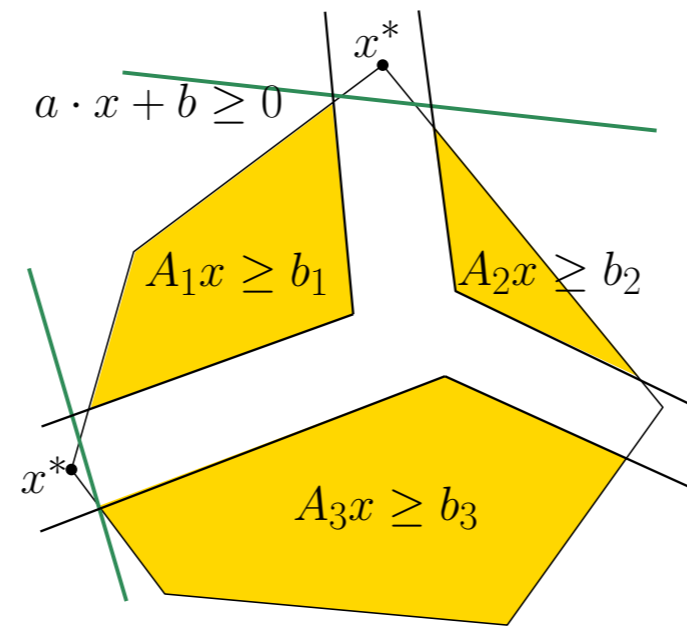


Cut Generating SDPs : A disjunctive
framework for MINLPs
with polynomial constraints

Amitabh Basu and Sam Burer

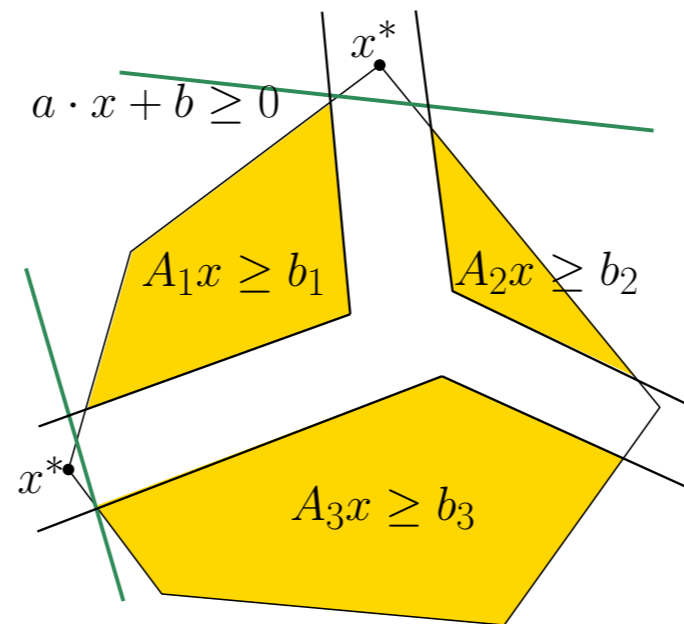
MINLP workshop, CMU Pittsburgh, June 2, 2014

Balas' CGLPs : A disjunctive framework for the polyhedral world



Balas' CGLPs : A disjunctive framework for the polyhedral world

Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

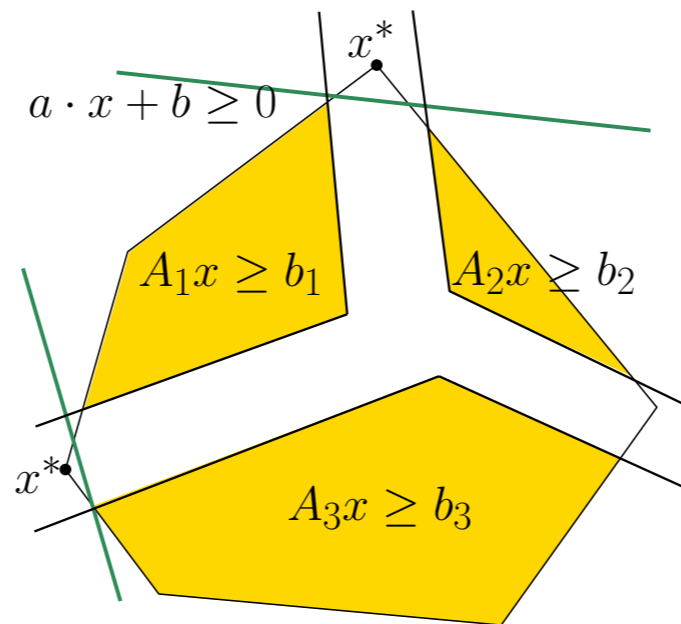
+ Normalizing constraints

Variables : a, b, λ_i

CGLP

Balas' CGLPs : A disjunctive framework for the polyhedral world

Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

Dual



$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

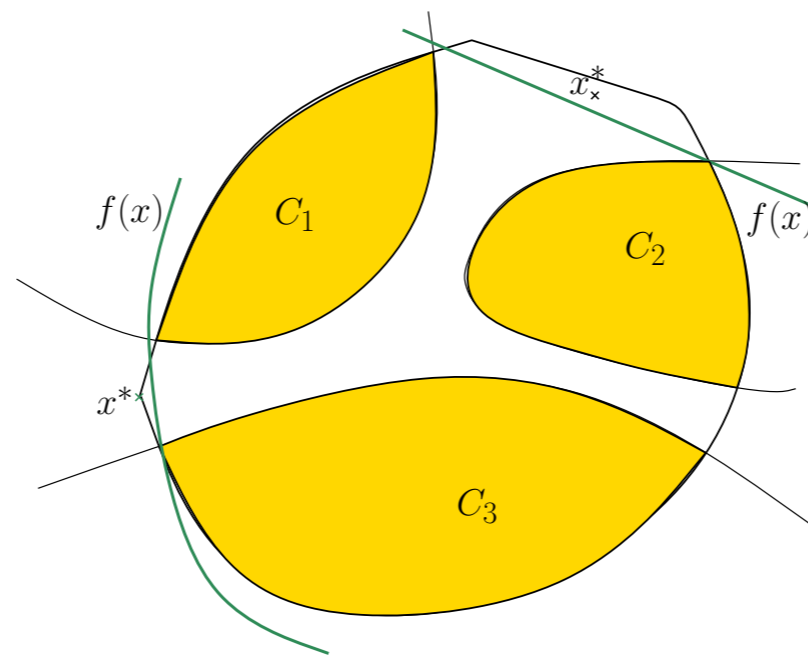
$$t_i \geq 0 \quad i = 1, 2, 3$$

CGLP

Extended Formulation

Cut Generating SDPs (CGSDPs) : A disjunctive framework for the polynomial constraints

Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

$$t_i \geq 0 \quad i = 1, 2, 3$$

Dual

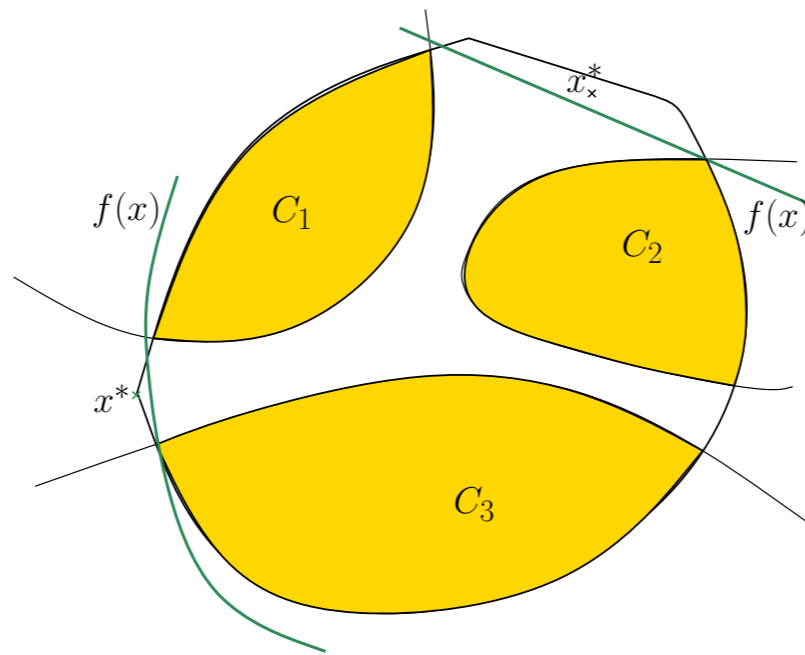


CGLP

Extended Formulation

Cut Generating SDPs (CGSDPs) : A disjunctive framework for the polynomial constraints

Positivstellensatz:
Nonlinear
Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

$$t_i \geq 0 \quad i = 1, 2, 3$$

Dual

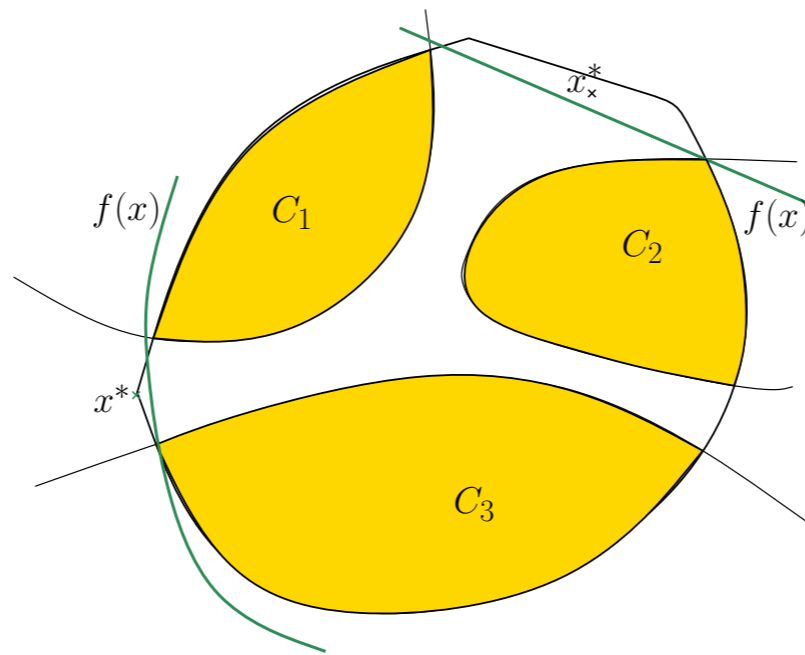


CGLP

Extended
Formulation

Cut Generating SDPs (CGSDPs) : A disjunctive framework for the polynomial constraints

Positivstellensatz:
Nonlinear
Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

Dual



$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

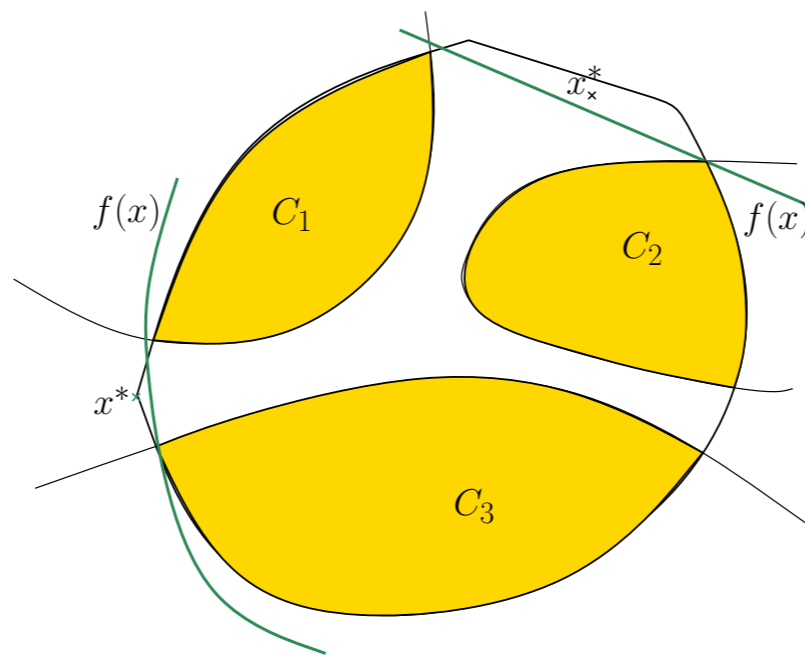
$$t_i \geq 0 \quad i = 1, 2, 3$$

CGSDP

Extended
Formulation

Cut Generating SDPs (CGSDPs) : A disjunctive framework for the polynomial constraints

Positivstellensatz:
Nonlinear
Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

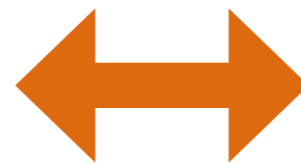
$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

Dual



$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

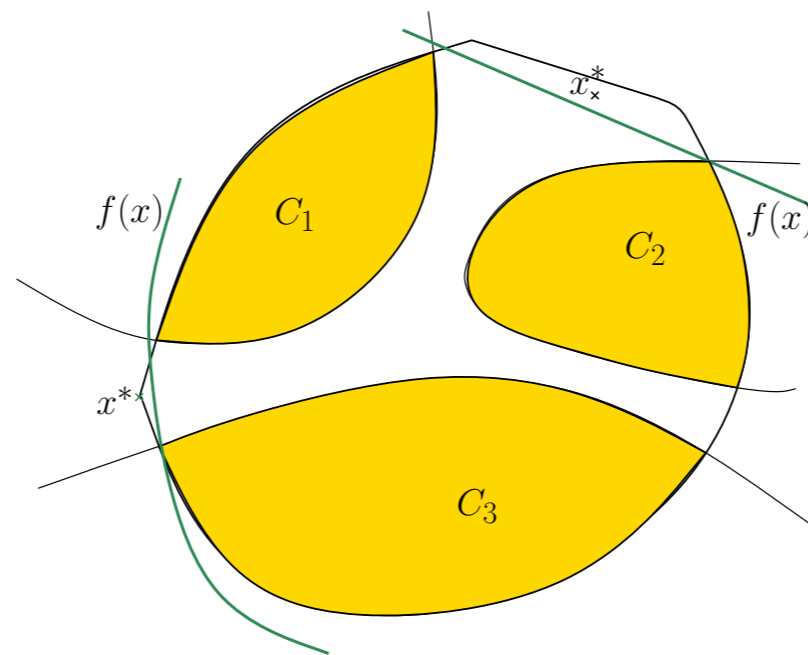
$$t_i \geq 0 \quad i = 1, 2, 3$$

CGSDP

Extended
Formulation

Cut Generating SDPs (CGSDPs) : A disjunctive framework for the polynomial constraints

Positivstellensatz:
Nonlinear
Farkas' Lemma



$$\min x^* \cdot a + b$$

$$\text{s.t. } a^T = \lambda_i^T A_i \quad i = 1, 2, 3$$

$$\lambda_i^T b_i + b \geq 0 \quad i = 1, 2, 3$$

$$\lambda_i \geq 0 \quad i = 1, 2, 3$$

+ Normalizing constraints

Variables : a, b, λ_i

SDP

Dual



$$A_i x_i - b_i t_i = 0 \quad i = 1, 2, 3$$

$$x_1 + x_2 + x_3 = x$$

$$t_1 + t_2 + t_3 = 1$$

$$t_i \geq 0 \quad i = 1, 2, 3$$

CGSDP

Extended
Formulation as
SDP

*See you at
my poster !!*